

Homework 4 for Columbia B9136

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Exercise 1. In this problem you are asked to generalize the RANKING algorithm and competitive ratio $1 - 1/e$ for online bipartite matching to the *vertex-weighted* setting, defined as follows. Each item i has a fixed, known reward $r_i > 0$ for getting matched (it can only get matched once). The objective is to maximize the total reward of the final matching.

- a) Write the new primal and dual LP's. (This requires a very minor modification.)
- b) The generalized RANKING algorithm no longer has a simple random-order interpretation, but is still described by drawing a random seed Y_i uniformly from $[0, 1]$ for each item i . The algorithm is now described by: for each t , match the available neighbor i with highest value of $r_i(1 - p(Y_i))$, where $p(y) = e^{y-1}$ is the same function as before. Small values of Y_i still boost item i 's priority. You should now think of item i as generating revenue $r_i p(Y_i)$ when sold, with the utility to the buying agent being $r_i(1 - p(Y_i))$.

Given this description of the algorithm, define dual variables whose objective is at most the algorithm's value divided by c , where $c = 1 - 1/e$.

- c) Argue that the constructed dual solution is feasible in expectation. The paradigm is the same, where we consider edge (i, t) and fix an arbitrary Y_{-i} . Find a generalized way of defining *threshold value* τ so that the following properties hold:
 - if $Y_i < \tau$, then i gets matched sooner or later;
 - β_t will be set to a value at least $r_i(1 - p(\tau))/c$.

(This is the hardest step, but still follows from the same argument as before, where we have to now be careful that β_t could be set based on an $i' \neq i$ with $r_{i'} \neq r_i$. Also, one new scenario that could arise is $\tau = 0$, i.e. even if $Y_i = 0$ item i is not guaranteed to get matched.)

- d) Finish the proof that generalized RANKING is $(1 - 1/e)$ -competitive under vertex weights. You can check your steps by setting $r_1 = \dots = r_n = 1$ and seeing if your proof coincides with the old proof.

Exercise 2. Find a counterexample to “Claim 2” from the analysis of RANKING, in the case of Adwords. In particular, consider the Greedy algorithm, and consider its execution on an instance \mathcal{I} and its parallel execution on an instance \mathcal{I}' that is identical to \mathcal{I} except one advertiser is removed. There is less competition from the advertiser side. Nonetheless, show that for an impression t , it is possible that the revenue earned from it is strictly higher in the parallel universe.